

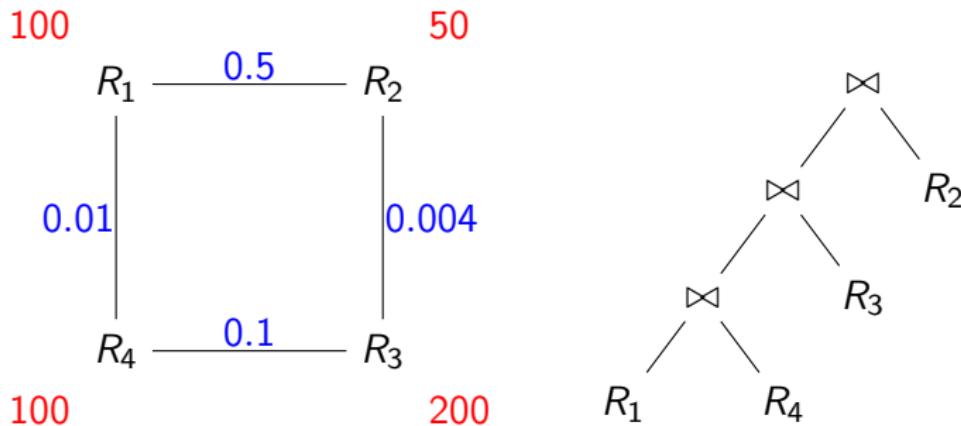
Query Optimization

Exercise Session 9

Andrey Gubichev

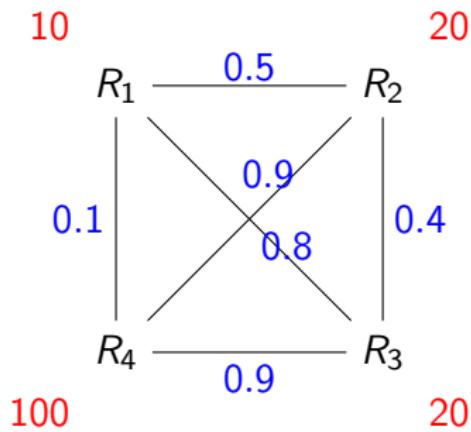
December 15, 2014

II example - 1 (Bui Nhat Nam, Nguyen Dinh Duy)



- Candidates: 3142, 4132, 3142, 4312, 3412

II example - 2



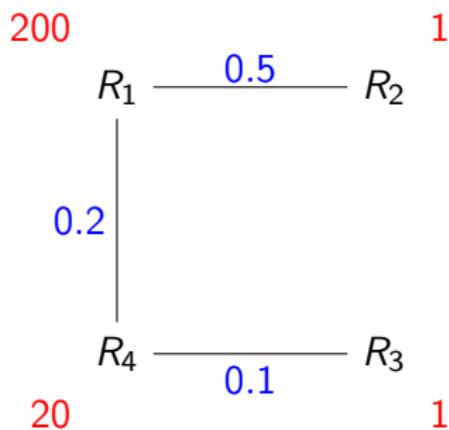
Start: $((R_1 R_2) R_3) R_4$: 6004

- ▶ assoc: $(R_1 (R_2 R_3)) R_4$: 6064
- ▶ assoc: $(R_1 R_2) (R_3 R_4)$: 7084
- ▶ l j ex: $((R_1 R_3) R_2) R_4$: 6084
- ▶ l j ex: $((R_1 R_2) R_4) R_3$: 6084

Optimal: $(R_1 R_4)(R_2 R_3)$: 5444

Order Preserving Joins: Example

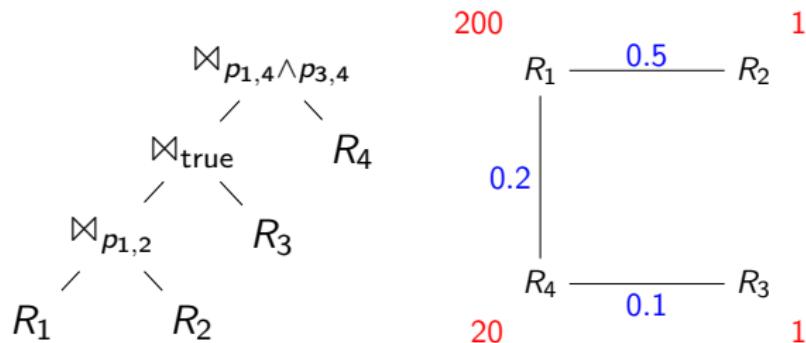
Consider the following sequence of relations R_1, R_2, R_3, R_4 and their join graph:



Give a fully-parenthesized, optimal join-expression that abides by this order. Use C_{out} as a cost function.

Order Preserving Joins: Baseline

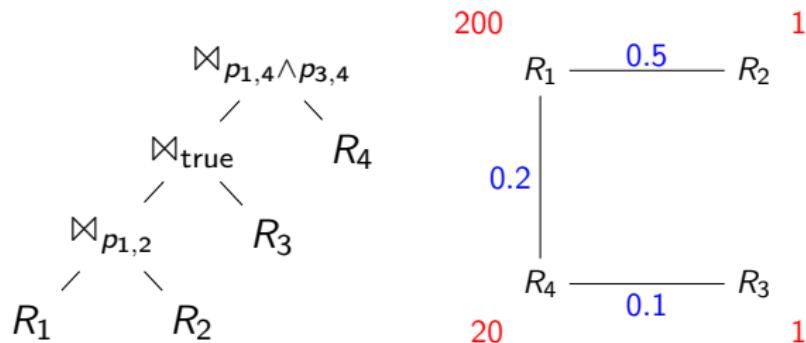
Let's start off with a cost analysis of the left-deep tree:



$$C_{out} =$$

Order Preserving Joins: Baseline

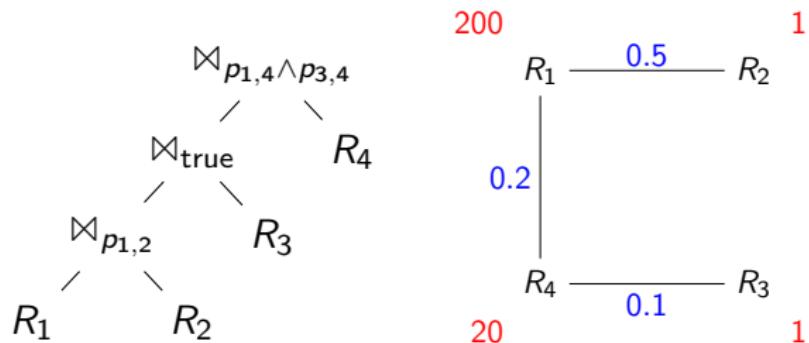
Let's start off with a cost analysis of the left-deep tree:



$$C_{out} = 100$$

Order Preserving Joins: Baseline

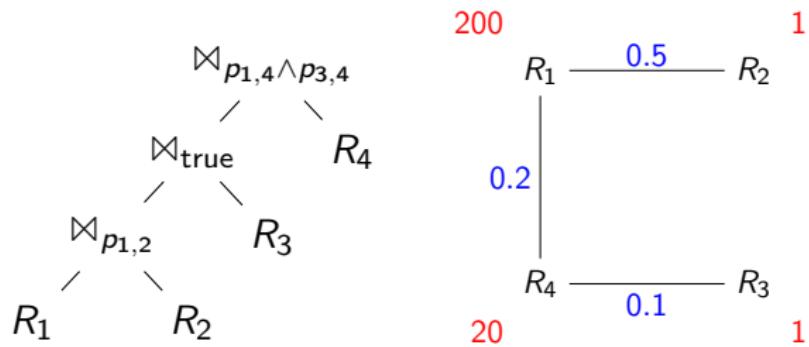
Let's start off with a cost analysis of the left-deep tree:



$$C_{out} = 100 + 100$$

Order Preserving Joins: Baseline

Let's start off with a cost analysis of the left-deep tree:



$$C_{out} = 100 + 100 + 40 = 240$$

Order Preserving Joins: Initialization

OrderPreservingJoins($R = \{R_1, \dots, R_n\}, P$)

Input: a set of relations to be joined and a set of predicates

Output: fills p, s, c, t

for each $1 \leq i \leq n$ {

$p[i, i]$ = predicates from P applicable to R_i

$P = P \setminus p[i, i]$

$s[i, i]$ = statistics for $\sigma_{p[i, i]}(R_i)$

$c[i, i]$ = costs for $\sigma_{p[i, i]}(R_i)$

}

predicates p		statistics s			costs c		
\emptyset					0		
	\emptyset			1		0	
		\emptyset			1		
			\emptyset			0	
					20		0

Order Preserving Joins: Constructing the Bushy Tree

```
01 for each  $2 \leq l \leq 4$  ascending (in text:  $2 \leq l \leq n$ )
02   for each  $1 \leq i \leq 5 - l$  (in text:  $1 \leq i \leq n - l + 1$ )
03      $j = i + l - 1$ 
04      $p[i, j] =$  predicates from  $P$  applicable to  $R_i, \dots, R_j$ 
05      $P = P \setminus p[i, j]$ 
06      $s[i, j] =$  statistics derived from  $s[i, j - 1]$  and  $s[j, j]$  including  $p[i, j]$ 
07      $c[i, j] = \infty$ 
08   for each  $i \leq k < j$ 
10      $q = c[i, k] + c[k + 1, j] +$  costs for  $s[i, k]$  and  $s[k + 1, j]$  and  $p[i, j]$ 
11     if  $q < c[i, j]$ 
12        $c[i, j] = q$ 
13        $t[i, j] = k$ 
```

predicates p		statistics s			costs c			split points t		
\emptyset					0					
	\emptyset					0				
		\emptyset					0			
			\emptyset					0		

line =

$l =$

$i =$

$j =$

$k =$

$q =$

Order Preserving Joins: Constructing the Bushy Tree

```
01 for each  $2 \leq l \leq 4$  ascending (in text:  $2 \leq l \leq n$ )
02   for each  $1 \leq i \leq 5 - l$  (in text:  $1 \leq i \leq n - l + 1$ )
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11     if  $q < c[i, j]$ 
12        $c[i, j] = q$ 
13        $t[i, j] = k$ 
```

predicates p		statistics s			costs c		split points t	
\emptyset	{ $p_{1,2}$ }				0	∞		
	\emptyset				0	0		
		\emptyset				0		
			\emptyset			0		

line = 08

$I = 2$

$i = 1$

$j = 2$

$k =$

$q =$

Order Preserving Joins: Constructing the Bushy Tree

```
01 for each  $2 \leq l \leq 4$  ascending (in text:  $2 \leq l \leq n$ )
02   for each  $1 \leq i \leq 5 - l$  (in text:  $1 \leq i \leq n - l + 1$ )
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05      $P = P \setminus p[i, j]$ 
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11     if  $q < c[i, j]$ 
12        $c[i, j] = q$ 
13        $t[i, j] = k$ 
```

predicates p		statistics s				costs c		split points t		
\emptyset	{ $p_{1,2}$ }					0	100			
	\emptyset						0			
		\emptyset						0		
				\emptyset						

line = 13

$l = 2$

$i = 1$

$j = 2$

$k = 1$

$$q = 0 + 0 + 200 \cdot 1 \cdot \frac{1}{2} = 100$$

Order Preserving Joins: Constructing the Bushy Tree

```
01 for each  $2 \leq l \leq 4$  ascending (in text:  $2 \leq l \leq n$ )
02   for each  $1 \leq i \leq 5 - l$  (in text:  $1 \leq i \leq n - l + 1$ )
03      $j = i + l - 1$ 
04      $p[i, j] =$  predicates from  $P$  applicable to  $R_i, \dots, R_j$ 
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11     if  $q < c[i, j]$ 
12        $c[i, j] = q$ 
13        $t[i, j] = k$ 
```

predicates p		statistics s				costs c			split points t		
\emptyset	$\{p_{1,2}\}$					0	100				
	\emptyset	\emptyset					1	$\textcolor{blue}{1}$			
		\emptyset						1			
					200			20			

line = 11
 $I = 2$
 $i = 2$
 $j = 3$
 $k = 2$
 $q = 0 + 0 + 1 \cdot 1 \cdot 1 = 1$

Order Preserving Joins: Constructing the Bushy Tree

```
01 for each  $2 \leq l \leq 4$  ascending (in text:  $2 \leq l \leq n$ )
02   for each  $1 \leq i \leq 5 - l$  (in text:  $1 \leq i \leq n - l + 1$ )
03      $j = i + l - 1$ 
04      $p[i, j] =$  predicates from  $P$  applicable to  $R_i, \dots, R_j$ 
05      $P = P \setminus p[i, j]$ 
06      $s[i, j] =$  statistics derived from  $s[i, j - 1]$  and  $s[j, j]$  including  $p[i, j]$ 
07      $c[i, j] = \infty$ 
08   for each  $i \leq k < j$ 
10      $q = c[i, k] + c[k + 1, j] +$  costs for  $s[i, k]$  and  $s[k + 1, j]$  and  $p[i, j]$ 
11     if  $q < c[i, j]$ 
12        $c[i, j] = q$ 
13        $t[i, j] = k$ 
```

predicates p		statistics s				costs c			split points t		
\emptyset	$\{p_{1,2}\}$	200	100	1	20	0	100	0	1	2	
\emptyset	\emptyset										
	\emptyset			1							
					1						
						20					
								0			
									1		
										2	

line = 13
 $l = 2$
 $i = 2$
 $j = 3$
 $k = 2$
 $q = 1$

Order Preserving Joins: Constructing the Bushy Tree

```
01 for each  $2 \leq l \leq 4$  ascending (in text:  $2 \leq l \leq n$ )
02   for each  $1 \leq i \leq 5 - l$  (in text:  $1 \leq i \leq n - l + 1$ )
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06      $s[i, j] =$  statistics derived from  $s[i, j - 1]$  and  $s[j, j]$  including  $p[i, j]$ 
07      $c[i, j] = \infty$ 
08   for each  $i \leq k < j$ 
10      $q = c[i, k] + c[k + 1, j] +$  costs for  $s[i, k]$  and  $s[k + 1, j]$  and  $p[i, j]$ 
11     if  $q < c[i, j]$ 
12        $c[i, j] = q$ 
13        $t[i, j] = k$ 
```

predicates p		statistics s				costs c			split points t	
\emptyset	$\{p_{1,2}\}$					0	100			1
	\emptyset	\emptyset					1	1		
		\emptyset	$\{p_{3,4}\}$					1	2	
				\emptyset				20		

line = 11
 $I = 2$
 $i = 3$
 $j = 4$
 $k = 3$
 $q = 0 + 0 + 1 \cdot 20 \cdot \frac{1}{10} = 2$

Order Preserving Joins: Constructing the Bushy Tree

```
01 for each  $2 \leq l \leq 4$  ascending (in text:  $2 \leq l \leq n$ )
02   for each  $1 \leq i \leq 5 - l$  (in text:  $1 \leq i \leq n - l + 1$ )
03      $j = i + l - 1$ 
04      $p[i, j] =$  predicates from  $P$  applicable to  $R_i, \dots, R_j$ 
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11     if  $q < c[i, j]$ 
12        $c[i, j] = q$ 
13        $t[i, j] = k$ 
```

predicates p		statistics s				costs c			split points t		
\emptyset	$\{p_{1,2}\}$	200	100	1	1	0	100	0	1	2	3
\emptyset	$\{p_{1,2}\}$			1	1			0	1		
	\emptyset				1	2			0	2	
					20					0	

line = 13

$I = 2$

$i = 3$

$j = 4$

$k = 3$

$q = 2$

Order Preserving Joins: Calling extract-plan

i/j	1	2	3	4
1		1	1	1
2			2	3
3				3
4				

The values of t are:

$\text{ExtractPlan}(R = \{R_1, \dots, R_n\}, t, p)$

Input: a set of relations, arrays t and p

Output: a bushy join tree

return $\text{ExtractPlanRec}(R, t, p, 1, n)$

$\text{ExtractPlanRec}(R = \{R_1, \dots, R_n\}, t, p, i, j)$

if $i < j$

$T_1 = \text{ExtractPlanRec}(R, t, p, i, t[i, j])$

$T_2 = \text{ExtractPlanRec}(R, t, p, t[i, j] + 1, j)$

return $T_1 \bowtie_{p[i, j]}^L T_2$

else

return $\sigma_{p[i, j]} R_i$

Order Preserving Joins: extract-plan callstack

```
extract-subplan(..., i=1, j=4)
    extract-subplan(..., i=1, j=1)
    extract-subplan(..., i=2, j=4)
        extract-subplan(..., i=2, j=3)
            extract-subplan(..., i=2, j=2)
            extract-subplan(..., i=3, j=3)
    return (R2 ⋈true R3)
    extract-subplan(..., i=4, j=4)
return ((R2 ⋈true R3) ⋈p3,4 R4)
return (R1 ⋈p1,2 ∧ p1,4 ((R2 ⋈true R3) ⋈p3,4 R4))
```

The total cost of this plan is $c[1, 4] = 43$.

Info

- ▶ Submit exercises to Andrey.Gubichev@in.tum.de
- ▶ Due December 22, 2014.